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A number of laws are derived which establish relationships between throughput, response time, device utilization, space-time products and various other factors related to computer system performance. These laws are obtained by using the operational method of computer system analysis. The operational method, which differs significantly from the conventional stochastic modeling approach, is based on a set of concepts that correspond naturally and directly to observed properties of real computer systems. Operational methods simply refer to methods of physical observation and measurement, in other words, they are directly measurable or directly measured. Except for measurement errors, the operational laws apply with complete precision to all collections of observational data, and they are similar to fundamental laws found in other areas of engineering and applied science.

The three main performance evaluation techniques are:

1. Performance measurement

2. Analytic performance modelling

3. Simulation performance modelling

The performance measurement technique

On-chip Performance Monitoring Counters: All state-of-the-art high performance microprocessors including Intel's Pentium III and Pentium IV, IBM's POWER 3 and POWER 4 processors, AMD's Athlon, Compaq's Alpha, and Sun's UltraSPARC processors incorporate on-chip performance monitoring counters which can be used to understand performance of these microprocessors while they run complex, real-world workloads. This ability has overcome a serious limitation of simulators, that they often could not execute complex workloads. Now, complex run time systems involving multiple software applications can be evaluated and monitored very closely. All microprocessor vendors nowadays release information on their performance monitoring counters, although they are not part of the architecture. For illustration of on-chip performance monitoring, we use the Intel Pentium processors. The microprocessors in the Intel Pentium contain two performance monitoring counters.

Off-chip hardware measurement: Instrumentation using hardware means can also be done by attaching off-chip hardware..

Speed-Tracer from AMD: AMD developed this hardware tracing platform to aid in the design of their x86 microprocessors. When an application is being traced, the tracer interrupts the processor on each instruction boundary. The state of the CPU is captured on each interrupt and then transferred to a separate control machine where the trace is stored. The trace contains virtually all valuable pieces of information for each instruction that executes on the processor. Operating system activity can also be traced. However, tracing in this manner can be invasive, and may slow down the processor. Although the processor is running slower, external events such as disk and memory accesses still happen in real time, thus looking very fast to the slowed-down processor. Usually this issue is addressed by adjusting the timer interrupt frequency. Use of this performance monitoring facility can be seen in Merten and Bhargava.

**L**ogic Analyzers: Poursepanj and Christie use a Tektronix TLA 700 logic analyzer to analyze 3D graphics workloads on AMD-K6-2 based systems. Detailed logic analyzer traces are limited by restrictions on sizes and are typically used for the most important sections of the program under analysis. Preliminary coarse level analysis can be done by performance monitoring counters and software instrumentation. Poursepanj and Christie used logic analyzer traces for a few tens of frames which covered a second or two of smooth motion.

Software Monitoring: Software monitoring is often performed by utilizing architectural features such as a trap instruction or a breakpoint instruction on an actual system, or on a prototype. The VAX processor from Digital (now Compaq) had a T-bit that caused an exception after every instruction. Software monitoring used to be an important mode of performance evaluation before the advent of on-chip performance monitoring counters. The primary advantage of software monitoring is that it is easy to do. However, disadvantages include that the instrumentation can slow down the application. The overhead of servicing the exception, switching to a data collection process, and performing the necessary tracing can slow down a program by more than 1000 times. Another disadvantage is that software monitoring systems typically only handle the user activity.

Micro coded Instrumentation: Digital (now Compaq) used microcoded instrumentation to obtain traces of VAX and Alpha architectures. The ATUM tool [14] used extensively by Digital in the late 1980s and early 1990s uses microcoded instrumentation. This is a technique lying between trapping information on each instruction using hardware interrupts (traps) or software traps. The tracing system essentially modified the VAX microcode to record all instruction and data references in a reserved portion of memory. Unlike software monitoring, ATUM could trace all processes including the operating system. However, this kind of tracing is invasive, and can slow down the system by a factor of 10 without including the time to write the trace to the disk.

The Analytical performance techniques

Using the Analytical performance modelling, we can achieve an accurate performance report: Performance measurement as described in the previous section can be done only if the actual system or a prototype exists. It is expensive to build prototypes for early stage evaluation. Hence one needs to resort to some kind of modeling in order to study systems yet to be built. Performance modeling can be done using simulation models or analytical models.

The Simulation performance techniques

Trace-driven simulation: consists of a simulator model whose input is modeled as a trace or sequence of information representing the instruction sequence that would have actually executed on the target machine. A simple trace driven cache simulator needs a trace consisting of address values. Depending on whether the simulator is modeling a unified instruction or data cache, the address trace should contain addresses of instruction and data references. Cachesim5 and Dinero IV are examples of cache simulators for memory reference traces. These simulators are not timing simulators. There is no notion of simulated time or cycles, only references. They are not functional simulators. Data and instructions do not move in and out of the caches. The primary result of simulation is hit and miss information. The basic idea is to simulate a memory hierarchy consisting of various caches. The various parameters of each cache can be set separately (architecture, mapping policies, replacement policies, write policy, statistics). During initialization, the configuration to be simulated is built up, one cache at a time, starting with each memory as a special case.

Execution Driven Simulation: There are two meanings in which this term is used by researchers and practitioners. Some refer to simulators that take program executables as input as execution driven simulators. These simulators utilize the actual input executable and not a trace. Hence the size of the input is proportional to the static instruction count and not the dynamic instruction count. Mis-predicted branches can be accurately simulated as well. Thus, these simulators solve the two major problems faced by trace-driven simulators. The widely used Simple-scalar simulator is an example of such an execution driven simulator. With this tool set, the user can simulate real programs on a range of modern processors and systems, using fast execution-driven simulation.

Complete system simulation: Many execution and trace driven simulators only simulate the processor and memory subsystem. Neither I/O activity nor operating system activity is handled in simulators like Simple scalar. But in many 9 workloads, it is extremely important to consider I/O and operating system activity. Complete system simulators are complete simulation environments that model hardware components with enough detail to boot and run a full-blown commercial operating system. The functionality of the processors, memory subsystem, disks, buses, SCSI/IDE/FC controllers, network controllers, graphics controllers, CD-ROM, serial devices, timers, etc are modeled accurately in order to achieve this. While functionality stays the same, different microarchitectures in the processing component can lead to different performance. Most of the complete system simulators use micro-architectural models that can be plugged in and out. For instance, SimOS, a popular complete system simulator provides a simple pipelined processor model and an aggressive superscalar processor model. SimOS and SIMICS can simulate uniprocessor and multiprocessor systems.

Stochastic Discrete Event Driven Simulation: It is possible to simulate systems in such a way that the input is derived stochastically rather than as a trace/executable from an actual execution. For instance, one can construct a memory system simulator in which the inputs are assumed to arrive according to a Gaussian distribution. Such models can be written in general purpose languages such as C, or using special simulation languages such as SIMSCRIPT. Languages such as SIMSCRIPT have several built-in primitives to allow quick simulation of most kinds of common systems. There are built-in input profiles, resource templates, process templates, queue structures, etc. to facilitate easy simulation of common systems.

Program Profilers: There are a class of tools called software profiling tools, which are similar to simulators and performance measurement tools. These tools are used to generate traces, to obtain instruction mix, and a variety of instruction statistics. They can be thought of as software monitoring on a simulator. They input an executable and decode and analyze each instruction in the executable. These program profilers can be used as the front end of simulators.